

A NOTE ON COMPUTING THE GENERALIZED INVERSE $A_{T,S}^{(2)}$ OF A MATRIX A

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The generalized inverse $A_{T,S}^{(2)}$ of a matrix A is a $\{2\}$ -inverse of A with the prescribed range T and null space S . A representation for the generalized inverse $A_{T,S}^{(2)}$ has been recently developed with the condition $\sigma(GA|_T) \subset (0, \infty)$, where G is a matrix with $R(G) = T$ and $N(G) = S$. In this note, we remove the above condition. Three types of iterative methods for $A_{T,S}^{(2)}$ are presented if $\sigma(GA|_T)$ is a subset of the open right half-plane and they are extensions of existing computational procedures of $A_{T,S}^{(2)}$, including special cases such as the weighted Moore-Penrose inverse $A_{M,N}^\dagger$ and the Drazin inverse A^D . Numerical examples are given to illustrate our results.

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1. Introduction. Given a complex matrix $A \in \mathbb{C}^{m \times n}$, any matrix $X \in \mathbb{C}^{n \times m}$ satisfying $XAX = X$ is called a $\{2\}$ -inverse of A . Let T and S be subspaces of \mathbb{C}^n and \mathbb{C}^m , respectively. A matrix $X \in \mathbb{C}^{n \times m}$ is called a $\{2\}$ -inverse of A with the prescribed range T and null space S , denoted by $A_{T,S}^{(2)}$, if the following conditions are satisfied:

$$XAX = X, \quad R(X) = T, \quad N(X) = S, \quad (1.1)$$

where $R(X)$ is the range of X and $N(X)$ is the null space of X . It is a well-known fact [1] that if $\dim T = \dim S^\perp \leq \text{rank}(A)$, then there exists a unique $A_{T,S}^{(2)}$ if and only if $AT \oplus S = \mathbb{C}^m$. It is obvious from the definition above that $AA_{T,S}^{(2)} = P_{AT,S}$ and $A_{T,S}^{(2)}A = P_{T,(A^*S^\perp)^\perp}$, where P_{S_1,S_2} is the projector on the subspace S_1 along the subspace S_2 .

There are seven types of important $\{2\}$ -inverses of A : the Moore-Penrose inverse A^\dagger , the weighted Moore-Penrose inverse $A_{M,N}^\dagger$, the W -weighed Drazin inverse $A_{d,w}$, the Drazin inverse A^D , the group inverse $A^\#$, the Bott-Duffin inverse $A_{(L)}^{(-1)}$, and the generalized Bott-Duffin inverse $A_{(L)}^{(\dagger)}$. All of them are the special cases of the generalized inverse $A_{T,S}^{(2)}$ of A for specific T and S .

LEMMA 1.1. (a) Let $A \in \mathbb{C}^{m \times n}$ [1]. Then, for the Moore-Penrose inverse A^\dagger and the weighted Moore-Penrose inverse A_{MN}^\dagger ,

- (i) $A^\dagger = A_{R(A^*),N(A^*)}^{(2)}$;
- (ii) $A_{M,N}^\dagger = A_{R(N^{-1}A^*M),N(N^{-1}A^*M)}^{(2)}$, where N and M are Hermitian positive definite matrices of order n and m , respectively;
- (iii) $A_{d,w} = (WAW)^{(2)}_{R(A(WA)^q),N(A(WA)^q)}$, where $W \in \mathbb{C}^{n \times m}$ and $q = \text{Ind}(WA)$, the index of WA .

(b) [1, 2, 3] Let $A \in \mathbb{C}^{n \times n}$. Then, for the Drazin inverse A^D , the group inverse $A^\#$, the Bott-Duffin inverse $A_{(L)}^{(-1)}$, and the generalized Bott-Duffin inverse $A_{(L)}^{(\dagger)}$,

- (iv) $A^D = A_{R(A^k), N(A^k)}^{(2)}$, where $k = \text{Ind}(A)$;
- (v) in particular, when $\text{Ind}(A) = 1$, $A^\# = A_{R(A), N(A)}^{(2)}$;
- (vi) $A_{(L)}^{(-1)} = P_L (AP_L + P_{L^\perp})^{-1} = A_{L, L^\perp}^{(2)}$, where L is a subspace of \mathbb{C}^n such that $AL \oplus L^\perp = \mathbb{C}^n$ and P_L is the orthogonal projector on L ;
- (vii) $A_{(L)}^{(\dagger)} = A_{(S)}^{(-1)} = A_{S, S^\perp}^{(2)}$, where $S = R(P_L A)$.

The $\{2\}$ -inverse has many applications, for example, the application in the iterative methods for solving nonlinear equations [1, 9] and the applications to statistics [6, 7]. In particular, $\{2\}$ -inverse plays an important role in stable approximations of ill-posed problems and in linear and nonlinear problems involving rank-deficient generalized inverse [8, 12]. In literature, researchers have proposed many numerical methods for computing $A_{T,S}^{(2)}$, see [2, 3, 11, 13, 15, 16, 18].

As usual, we denote the spectrum and the spectral radius of A by $\sigma(A)$ and $\rho(A)$, respectively. The notation $\|\cdot\|$ stands for the spectral norm. The following theorem applied in this note is from the theory of semi-iterative method.

THEOREM 1.2 (see [5]). *Let $B \in \mathbb{C}^{n \times n}$ be a nonsingular matrix and let $\sigma(B) \subset \Omega$, where Ω is a simply connected compact set excluding origin. If a sequence of polynomials $\{s_m(z)\}_{m=0}^\infty$ uniformly converges to $1/z$ on Ω , then $\{s_m(B)\}$ converges to B^{-1} .*

In this note, a representation for the generalized inverse $A_{T,S}^{(2)}$ with a condition $\sigma(GA|_T) \subset \{z : \text{Re}(z) > 0\}$, where G is a matrix with $R(G) = T$ and $N(G) = S$ is presented in Section 2. Euler-Knopp iterative method and semi-iterative methods for $A_{T,S}^{(2)}$ with linear convergence are derived in Section 3. Quadratically convergent methods for $A_{T,S}^{(2)}$ are developed in Section 4. Finally, numerical examples are given to illustrate our results.

2. Representation. In this section, we give a representation for the generalized inverse $A_{T,S}^{(2)}$, which may be viewed as an application of the classical theory summability to the representation of generalized inverse.

LEMMA 2.1 (see [13]). *Suppose $A \in \mathbb{C}^{m \times n}$. Let T and S be subspaces of \mathbb{C}^n and \mathbb{C}^m , respectively, such that $AT \oplus S = \mathbb{C}^m$. Suppose that $G \in \mathbb{C}^{n \times m}$ satisfies $R(G) = T$ and $N(G) = S$. Denote by $\tilde{A} = (GA)|_T$ the restriction of GA on T . Then $\text{Ind}(GA) = 1$ and*

$$A_{T,S}^{(2)} = \tilde{A}^{-1}G. \quad (2.1)$$

It follows from Lemma 1.1 that the existence of G is assured for each of the common seven types of generalized inverses: A^* , $N^{-1}A^*M$, $A(WA)^q$, A^k , A , P_L , and P_S . Now we are in a position to establish a presentation theorem.

THEOREM 2.2. *Let A , T , S , G , and \tilde{A} be as in Lemma 2.1. If $\sigma(\tilde{A})$ is contained in a simply connected compact set Ω excluding origin and a polynomial sequence $\{s_m(z)\}$ uniformly converges to $1/z$ on Ω , then*

$$A_{T,S}^{(2)} = \lim_{m \rightarrow \infty} s_m(\tilde{A})G. \quad (2.2)$$

Furthermore,

$$\frac{\|s_m(\tilde{A})G - A_{T,S}^{(2)}\|_P}{\|A_{T,S}^{(2)}\|_P} \leq \max_{z \in \sigma(\tilde{A})} |zs_m(z) - 1| + O(\epsilon), \quad (2.3)$$

where P is invertible such that $P^{-1}GAP$ is the ϵ -Jordan canonical form of GA and $\|B\|_P = \|P^{-1}B\|$ for each $B \in \mathbb{C}^{n \times m}$.

PROOF. Assume that $\sigma(\tilde{A}) \subset \Omega$. With applying [Theorem 1.2](#), we get

$$\lim_{m \rightarrow \infty} s_m(\tilde{A}) = \tilde{A}^{-1} \quad (2.4)$$

uniformly on Ω . It follows from [Lemma 2.1](#) that

$$\lim_{m \rightarrow \infty} s_m(\tilde{A})G = \tilde{A}^{-1}G = A_{T,S}^{(2)}. \quad (2.5)$$

The error can be written as

$$s_m(\tilde{A})G - A_{T,S}^{(2)} = (s_m(\tilde{A})\tilde{A} - I)A_{T,S}^{(2)}. \quad (2.6)$$

Since P is nonsingular such that $P^{-1}GAP$ is the ϵ -Jordan canonical form of GA , it is well known that

$$\|P^{-1}GAP\| \leq \rho(GA) + \epsilon. \quad (2.7)$$

Thus

$$\begin{aligned} \|s_m(\tilde{A})G - A_{T,S}^{(2)}\|_P &= \|(P^{-1}(s_m(\tilde{A})\tilde{A} - I)PP^{-1}A_{T,S}^{(2)})\|_P \\ &\leq \|P^{-1}(s_m(\tilde{A})\tilde{A} - I)P\| \|A_{T,S}^{(2)}\|_P \\ &\leq \left[\max_{z \in \sigma(\tilde{A})} |s_m(z)z - 1| + O(\epsilon) \right] \|A_{T,S}^{(2)}\|_P. \end{aligned} \quad (2.8)$$

The last inequality is based on the spectrum mapping since $s_m(z)$ is a polynomial in z . This completes the proof. \square

In order to make use of this general error estimate in [Theorem 2.2](#) on specific approximation procedures, it will be convenient to have lower and upper bounds for $\sigma(\tilde{A})$. This is given in the next lemma.

LEMMA 2.3. *Let A , T , S , G , and \tilde{A} be as in [Lemma 2.1](#). Then for each $\lambda \in \sigma(\tilde{A})$,*

$$\frac{1}{\|(GA)^\#\|} \leq |\lambda| \leq \|GA\|. \quad (2.9)$$

PROOF. We only show the first inequality since the second is trivial. It follows from [Lemma 2.1](#) that $\text{Ind}(GA) = 1$. Then the Jordan canonical form of GA is

$$GA = P \begin{bmatrix} C & 0 \\ 0 & 0 \end{bmatrix} P^{-1}, \quad (GA)^\# = P \begin{bmatrix} C^{-1} & 0 \\ 0 & 0 \end{bmatrix} P^{-1}, \quad (2.10)$$

where C is invertible. For each $\lambda \in \sigma(\tilde{A})$, $1/\lambda \in \sigma(\tilde{A}^{-1})$ since \tilde{A} is invertible. Consequently, we have

$$\frac{1}{|\lambda|} \leq \rho(\tilde{A}^{-1}) = \rho(C^{-1}) \leq \|(GA)^\#\|, \quad (2.11)$$

which leads to (2.9). This completes the proof. \square

REMARK 2.4. [Theorem 2.2](#) extends the representation of $A_{T,S}^{(2)}$ in [15] in which $\sigma(GA|_T) \subset (0, \infty)$ is required. The theorem also recovers the representations of A^D in [16] and $A_{M,N}^\dagger$ in [17] as special cases.

3. Iterative methods for $A_{T,S}^{(2)}$. In this section, we present applications of [Theorem 2.2](#) and [Lemma 2.3](#) in developing specific computational procedures for the generalized inverse $A_{T,S}^{(2)}$ and estimating corresponding error bounds.

A well-known summability method is called the Euler-Knopp method. A series $\sum_{m=0}^{\infty} a_m$ is said to be Euler-Knopp summable with parameter $\alpha > 0$ to the value a if the sequence defined by

$$s_m = \alpha \sum_{i=0}^m \sum_{j=0}^i \binom{i}{j} (1-\alpha)^{i-j} \alpha^j a_j \quad (3.1)$$

converges to a . If we choose $a_m = (1-z)^m$, $m \geq 0$, then as the Euler-Knopp transform of the series $\sum_{m=0}^{\infty} (1-z)^m$, we obtain a sequence $\{s_m(z)\}$, where

$$s_m(z) = \alpha \sum_{j=0}^m (1-\alpha z)^j. \quad (3.2)$$

Clearly, $\lim_{m \rightarrow \infty} s_m(z) = 1/z$ uniformly on any compact subset of an open set $E_\alpha := \{z : |1-\alpha z| < 1\}$. We assume that $\sigma(\tilde{A}) \subset \{z : \operatorname{Re}(z) > 0\}$. Denote

$$\phi := \max_{\lambda \in \sigma(\tilde{A})} \left\{ |\operatorname{Arg} \lambda| : -\frac{\pi}{2} < \operatorname{Arg} \lambda < \frac{\pi}{2} \right\}. \quad (3.3)$$

It follows from [Lemma 2.3](#) that

$$\sigma(\tilde{A}) \subset \{z = r e^{i\theta} : r_1 \leq r \leq r_2, -\phi \leq \theta \leq \phi\} =: F, \quad (3.4)$$

where $r_1 = 1/\|(GA)^\#\|$ and $r_2 = \|GA\|$. It can be shown with the law of Sines that

$$F \subset \{w : |w - g| \leq g\}, \quad \text{for } g = \frac{\|GA\|}{2 \cos \phi}. \quad (3.5)$$

If a parameter α satisfies

$$0 < \alpha < \frac{2 \cos \phi}{\|GA\|}, \quad (3.6)$$

then $\sigma(\tilde{A}) \subset E_\alpha$. There is always a simply connected compact set Ω such that $\sigma(\tilde{A}) \subset \Omega \subset E_\alpha$. Hence $s_m(z)$ of (3.2) uniformly converges to $1/z$ on Ω . It follows from [Theorem 2.2](#) that

$$A_{T,S}^{(2)} = \alpha \sum_{n=0}^{\infty} (I - \alpha G A)^n G. \quad (3.7)$$

Notice that if A_m is the m th partial sum, that is, $A_m = \alpha \sum_{j=0}^m (I - \alpha G A)^j G$, then an iteration form for $\{A_m\}$ is given by

$$A_0 = \alpha G, \quad A_{m+1} = (I - \alpha G A) A_m + \alpha G, \quad m \geq 0. \quad (3.8)$$

For an error bound, we note that the sequence of polynomials $\{s_m(z)\}$ satisfies

$$zs_{m+1}(z) - 1 = (1 - \alpha z)(zs_m(z) - 1). \quad (3.9)$$

Thus

$$|zs_m(z) - 1| = |1 - \alpha z|^{m+1} \leq \beta^{m+1} \rightarrow 0, \quad (m \rightarrow \infty), \quad (3.10)$$

where

$$\beta = \max_{z \in \sigma(\tilde{A})} |1 - \alpha z| \leq \max_{z \in F} |1 - \alpha z| < 1. \quad (3.11)$$

Actually, by the maximum modular theorem, $\max_{z \in F} |1 - \alpha z| = \max_{z \in \partial F} |1 - \alpha z|$. We denote four parts of ∂F as follows:

$$\begin{aligned} \Gamma_1 &= \{r_1 e^{i\theta} : -\phi \leq \theta \leq \phi\}, & \Gamma_2 &= \{r e^{i\phi} : r_1 \leq r \leq r_2\}, \\ \Gamma_3 &= \{r_2 e^{i\theta} : -\phi \leq \theta \leq \phi\}, & \Gamma_4 &= \{r e^{-i\phi} : r_1 \leq r \leq r_2\}. \end{aligned} \quad (3.12)$$

If $z \in \Gamma_1$, then $|1 - \alpha z|^2 = 1 - 2\alpha r_1 \cos \theta + \alpha^2 r_1^2$ and it is obvious that

$$\max_{z \in \Gamma_1} |1 - \alpha z| = |1 - \alpha r_1 e^{i\phi}|. \quad (3.13)$$

With an analogous argument, we have

$$\max_{z \in \Gamma_3} |1 - \alpha z| = |1 - \alpha r_2 e^{i\phi}|. \quad (3.14)$$

If $z \in \Gamma_2 \cup \Gamma_4$, then $|1 - \alpha z|^2 = 1 - 2\alpha r \cos \phi + \alpha^2 r^2$ is a quadratic function of r on $[r_1, r_2]$, which achieves its maximum at either $r = r_1$ or $r = r_2$. So

$$\max_{z \in \Gamma_2 \cup \Gamma_4} |1 - \alpha z| = \max \{|1 - \alpha r_1 e^{i\phi}|, |1 - \alpha r_2 e^{i\phi}|\}. \quad (3.15)$$

It follows from [Theorem 2.2](#) that an error bound is given by

$$\frac{\|A_m - A_{T,S}^{(2)}\|_P}{\|A_{T,S}^{(2)}\|_P} \leq \beta^{m+1} + O(\varepsilon), \quad (3.16)$$

where

$$\beta \leq \max \{ |1 - \alpha e^{i\phi} / \| (GA)^{\#} \| |, |1 - \alpha e^{i\phi} \| GA \| | \}. \quad (3.17)$$

Therefore, we have shown the following general convergence theorem.

THEOREM 3.1. *Let A, T, S , and G be as in Lemma 2.1. Suppose the spectrum of $GA|_T$ is contained in the open right half-plane. Then the sequence $\{A_m\}$ of (3.8) linearly converges to $A_{T,S}^{(2)}$, if $0 < \alpha < 2 \cos \phi / \|GA\|$, where ϕ is given by (3.3). Moreover, the relative error is bounded by (3.16).*

We remark that Theorem 3.1 is an extension of corresponding results in [15, 16].

The procedure of semi-iterative methods [5, 10] for solving a linear system can easily be extended to solve

$$X = HX + C, \quad \text{for } C \in \mathbb{C}^{n \times n}. \quad (3.18)$$

If $\rho(H) < 1$, then a sequence of matrices $\{X_m\}$, yielded by

$$X_0 = C; \quad X_{m+1} = HX_m + C \quad (m \geq 0), \quad (3.19)$$

converges to $(I - H)^{-1}C$. In general, let $1 \notin \sigma(H)$. As usual, based on a sequence of polynomials $\{p_m(z)\}$ given by

$$p_m(z) = \sum_{i=0}^m \pi_{m,i} z^i, \quad \text{where } \sum_{i=0}^m \pi_{m,i} = 1, \quad (3.20)$$

the corresponding semi-iterative method induced by $\{p_m(z)\}$ for the computation of $(I - H)^{-1}C$ is defined as

$$Y_m = \sum_{i=0}^m \pi_{m,i} X_i, \quad m \geq 0. \quad (3.21)$$

Moreover, the matrices Y_m and the corresponding residual matrices R_m are given by

$$Y_m = p_m(H)Y_0 + q_{m-1}(H)C, \quad R_m = p_m(H)(C - (I - H)Y_0), \quad (3.22)$$

where

$$q_{m-1}(z) = (1 - p_m(z)) / (1 - z) \quad \text{with } q_{-1}(z) = 0. \quad (3.23)$$

If $\{q_m(H)\}$ converges to $(I - H)^{-1}$, or equivalently, if $\{p_m(H)\}$ converges to 0, then the sequence $\{Y_m\}$ of (3.21) converges to $(I - H)^{-1}C$. Especially, for $H = I - GA|_T$ and $C = G$, $\{Y_m\}$ converges to $A_{T,S}^{(2)}$. With an application of Theorem 1.2, we have the following corollary.

COROLLARY 3.2. *Let A, T, S , and G be as in Lemma 2.1 and let $H = I - GA|_T$. If $\sigma(H)$ is contained in Ω_1 , a simply connected compact set excluding 1, and $\{q_m(z)\}$ of (3.23) uniformly converges to $1/(1 - z)$ on Ω_1 , then the sequence $\{Y_m\}$ of (3.21) converges to $A_{T,S}^{(2)}$ for $Y_0 = G$.*

Especially, Ω_1 is either a complex segment $[\alpha, \beta]$ excluding 1 or a closed ellipse in the left half-plane $\{z : \operatorname{Re}(z) < 1\}$ with foci α and β . Let a sequence of polynomials $\{p_m(z)\}$ given by

$$p_m(z) = \frac{T_m((z-\delta)/\xi)}{T_m((1-\delta)/\xi)}, \quad \left(\delta = \frac{\alpha+\beta}{2}, \quad \xi = \frac{\beta-\alpha}{2} \right), \quad (3.24)$$

where T_m is the m th Chebyshev polynomial. The semi-iterative method induced by $\{p_m(z)\}$ is the Chebyshev iterative method optimal for ellipse Ω_1 . The corresponding two-step stationary method with the same asymptotically optimal convergence rate is given by

$$\begin{aligned} Y_0 &= G; & Y_1 &= \mu(HY_0 + G); \\ Y_{m+1} &= \mu_0(HY_m + G) + \mu_1 Y_m + \mu_2 Y_{m-1}, & (m \geq 1), \end{aligned} \quad (3.25)$$

where

$$\mu_0 = \frac{4}{(\sqrt{1-\beta} + \sqrt{1-\alpha})^2}, \quad \mu_1 = \frac{\alpha+\beta}{2}\mu_0, \quad \mu_2 = 1 - \mu_0 - \mu_1. \quad (3.26)$$

The sequence $\{Y_m\}$ converges asymptotically optimally to $A_{T,S}^{(2)}$.

4. Quadratically convergent methods. Newton-Raphson method for finding the root $1/z$ of the function $s(w) = w^{-1} - z$ is given by

$$w_{m+1} = w_m(2 - zw_m), \quad \text{for a suitable } w_0. \quad (4.1)$$

For $\alpha > 0$, a sequence of functions $\{s_m(z)\}$ is defined by

$$s_0(z) = \alpha, \quad s_{m+1}(z) = s_m(z)[2 - zs_m(z)]. \quad (4.2)$$

Let $z \in \sigma(GA|_T)$ and $0 < \alpha < 2\cos\phi/\|GA\|$. It follows from the recursive form $zs_{m+1}(z) - 1 = -[zs_m(z) - 1]^2$ that

$$|zs_m(z) - 1| = |\alpha z - 1|^{2^m} \leq \beta^{2^m} \rightarrow 0, \quad \text{as } m \rightarrow \infty, \quad (4.3)$$

where an upper bound of β is given by (3.17).

The great attraction of the Newton-Raphson method is the generally quadratic nature of the convergence. Using the above facts in conjunction with Lemma 2.3, we see that a sequence $\{s_m(\tilde{A})\}$ defined by

$$s_0(\tilde{A}) = \alpha I, \quad s_{m+1}(\tilde{A}) = s_m(\tilde{A})[2I - \tilde{A}s_m(\tilde{A})] \quad (4.4)$$

has the property that $\lim_{m \rightarrow \infty} s_m(\tilde{A})G = A_{T,S}^{(2)}$. If we set $A_m = s_m(\tilde{A})G$, then

$$A_0 = \alpha G, \quad A_{m+1} = A_m(2I - AA_m). \quad (4.5)$$

Thus we have the following corollary.

COROLLARY 4.1. *Let A , T , S , and G be as in Lemma 2.1. Suppose that the spectrum of $\sigma(GA|_T) \subset \{z : \operatorname{Re}(z) > 0\}$. Then the sequence $\{A_m\}$ of (4.5) quadratically converges to $A_{T,S}^{(2)}$, for $0 < \alpha < 2 \cos \phi / \|GA\|$. Furthermore, an error bound is given by*

$$\|A_m - A_{T,S}^{(2)}\|_P \leq (\beta^{2^m} + O(\epsilon)) \|A_{T,S}^{(2)}\|_P, \quad (4.6)$$

where an upper bound of β is given in (3.17).

We remark that Corollary 4.1 is an extension of [4, 13, 15]. It covers iterative methods for $A_{M,N}^\dagger$ in [17].

The Newton-Raphson procedure can be speeded up by the successive matrix squaring technique in [14] if two parallel processors are available. In fact, the sequence in (4.5) is mathematically equivalent to

$$\begin{aligned} A_0 &= \alpha G, & P_0 &= I - \alpha G A, \\ A_{m+1} &= (I + P_m) A_m, & P_{m+1} &= P_m^2. \end{aligned} \quad (4.7)$$

There are two matrix multiplications each step both in (4.5) and (4.7). However, A_{m+1} and P_{m+1} in (4.7) can be calculated simultaneously.

Two algorithms given by (3.8) and (4.5) are also valid in the case when the spectrum of \tilde{A} is contained in the left half-plane with slight modification.

Moreover, all results in the previous two sections are valid without the restriction on $\sigma(GA)$ if G is substituted by another matrix. This is stated as the following corollary.

COROLLARY 4.2. *Let A , T , S , and G be as in Lemma 2.1. Then Theorem 3.1 and Corollaries 3.2 and 4.1 are valid without any restriction on the spectrum of $GA|_T$ if G is substituted by*

$$G_0 = G(GAG)^*G. \quad (4.8)$$

PROOF. It suffices to show that

$$R(G_0) = R(G), \quad N(G_0) = N(G), \quad \sigma(G_0 A|_T) \subset (0, \infty). \quad (4.9)$$

As a matter of fact (4.9) is a direct result of [4, Lemma 3.4]. \square

We remark a disadvantage of the choice G_0 of (4.8). In the case of computing A^D with $\operatorname{Ind}(A) = k \geq 3$, $G_0 = A^k (A^{2k+1})^* A^k$, the condition number of $G_0 A|_T$ will be extremely large since $\operatorname{cond}(G_0 A|_T) = \operatorname{cond}(A|_T)^{4k+2}$. An accurate numerical solution cannot be obtained if there is any round-off error in A .

5. Examples. Three examples are given in this section to illustrate the computations of three types of $A_{T,S}^{(2)}$. All calculations were performed on a PC with MATLAB.

EXAMPLE 5.1. Let A and W be 20 by 10 and 10 by 10 random matrices with entries on $[-1, 1]$, respectively. We choose M and N as random symmetric and positive definite matrices of order 20 and 10, respectively. The stop criterion in (4.5) is

$\|A_m - A_{m-1}\|_\infty \leq \epsilon = 10^{-10}$. Three special cases A^\dagger , $A_{M,N}^\dagger$, and $A_{d,w}$ are computed in this example. The choices of G , the number of iterations required and the norm of errors are listed in [Table 5.1](#).

TABLE 5.1. Newton-Raphson method for A^\dagger , $A_{M,N}^\dagger$, and $A_{d,w}$.

| $A_{T,S}^{(2)}$ | G | m | $\ A_m - A_{m-1}\ _\infty$ | $\ A_m - A_{T,S}^{(2)}\ _\infty$ |
|-------------------|-------------------|-----|----------------------------|----------------------------------|
| A^\dagger | A^* | 11 | 3.25E-14 | 2.56E-15 |
| $A_{M,N}^\dagger$ | $N^{-1}A^*M$ | 25 | 1.03E-15 | 3.09E-15 |
| $A_{d,w}$ | $AWA((WA)^*)^4WA$ | 36 | 1.96E-13 | 1.76E-07 |

It is remarked that the better accuracy of $A_{d,w}$ never be achieved and 1.7E-07 is the best error of $\|A_m - A_{d,w}\|_\infty$ even if $\|A_m - A_{m-1}\|_\infty \leq \epsilon = 10^{-10}$ is used as a stop criterion. This is because the condition number of $GWAW|_T$ is as large as 10^{10} . If 2-step semi-iterative method of (3.25) is applied to compute $A_{T,S}^{(2)}$, then $\{Y_m\}$ converges to A^\dagger after 54 iterations. However, the method fails to converge after 1500 iterations in other two cases because the segments $[\alpha, \beta]$ containing $\sigma(GA|_T)$ are $[-187970, 0.796]$ and $[-355800, 0.9997]$, respectively, so that the rate of asymptotic convergence is too slow.

EXAMPLE 5.2. Let A be 8 by 8 matrix with a complex spectrum given by

$$A = \begin{bmatrix} \frac{3}{2} & \frac{1}{3} & 0 & 0 & 0 & 0 & 0 & 0 \\ -\frac{1}{4} & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & -1 & \frac{3}{4} & -\frac{3}{4} & 0 & 0 & 0 & 0 \\ -1 & -1 & -\frac{3}{4} & \frac{3}{4} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & \frac{3}{4} & -\frac{3}{4} & -1 & -1 \\ 0 & 0 & -1 & 0 & -\frac{3}{4} & \frac{3}{4} & -1 & -1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & -\frac{1}{4} \\ 0 & 0 & 0 & 0 & 0 & 0 & \frac{1}{3} & \frac{3}{2} \end{bmatrix}. \quad (5.1)$$

In order to compute A^D , we choose $G = A^2$ since $\text{Ind}(A) = 2$. The spectrum of $GA|_T$, $\sigma(GA|_T) = \{1.875 \pm 0.674i, 1.875 \pm 0.674i, 3.375, 3.375\}$, is located on the right half-plane. The foci $\alpha = -2.3$ and $\beta = -0.5$ of an ellipse containing $\sigma(I - GA|_T)$ is selected. It requires 28 iterations of 2-step method of (3.25) to compute A^D with the ∞ -norm of the error less than 10^{-10} . As expected, Newton-Raphson algorithm of (4.5) converges much faster. It achieves the same accuracy with only 8 iterations.

TABLE 5.2. $A_{T,S}^{(2)}$ of a Toeplitz matrix.

| $A_{T,S}^{(2)}$ | G | No. of it. by Newton's | No. of it. by SIM |
|-------------------|-------------------|------------------------|-------------------|
| A^\dagger | A^* | 10 | 63 |
| $A_{M,N}^\dagger$ | $N^{-1}A^*M$ | 11 | 75 |
| $A_{d,w}$ | $AWA((WA)^*)^4WA$ | 31 | > 1500 |

EXAMPLE 5.3. Let r and c be a row vector and column vector, respectively, such that

$$\begin{aligned} r_1 &= c_1 = 2.5, \\ r_j &= \frac{(-1)^j j}{16} + \frac{i(j-1)}{j}, \quad \text{for } j = 2, 3, \dots, 16, \\ c_k &= \frac{(-1)^k k}{10}, \quad \text{for } k = 2, 3, \dots, 10. \end{aligned} \quad (5.2)$$

A 10×16 complex Toeplitz matrix A is constructed by r and c . The stop criterion is the same as in [Example 5.1](#). M and N are chosen positive definite diagonal matrix related to A , and W is a random matrix. The numbers of iterations by Newton's method and SIM method for $A_{T,S}^{(2)}$ are shown in [Table 5.2](#).

The data shows that Newton's method is much faster than that of SIM.

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